

# The Effectiveness of Subsidies and Taxes in Atomic Congestion Games

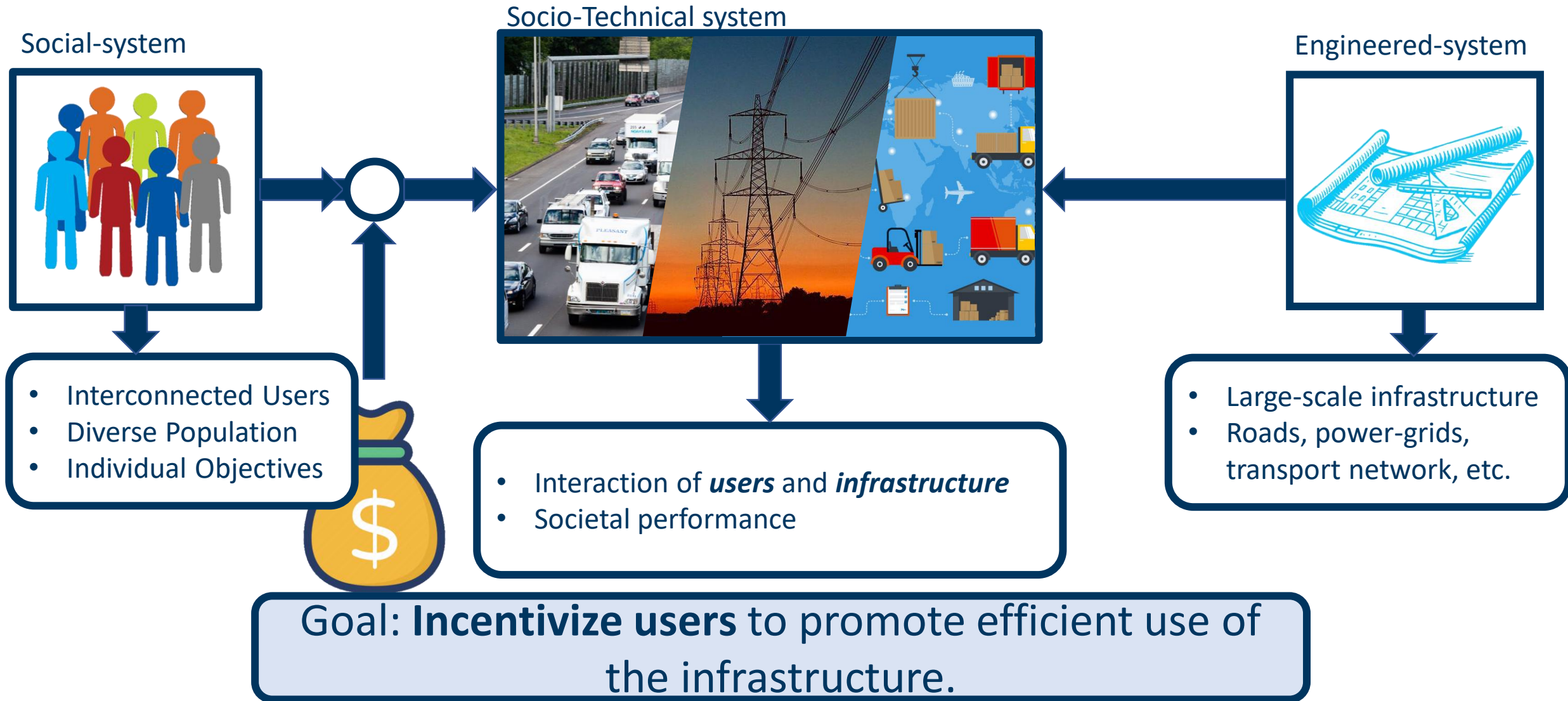


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# Socio-Technical Systems



# Carrots vs Sticks

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## Taxes/Tolls

(+)

Added cost to users



Tolls

A. De Palma, R. Lindsey, "Private roads, competition, and incentives to adopt time-based congestion tolling," *Elsevier*

Q. Wang, M. Liu, R. Jain, "Dynamic pricing of power in smart-grid networks," *IEEE Conference on Decision and Control*

M. Christopher, J. Gattorna, "Supply chain cost management and value-based pricing," *Elsevier*



Transportation

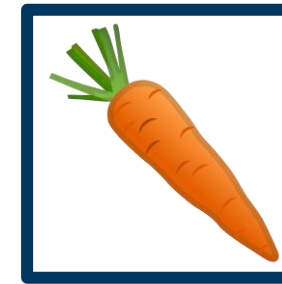
Power Grids

Supply-chain  
Management

## Subsidies/Rebates

(-)

Reduced cost to users



Subsidies

P. Maillé and N. E. Stier-Moses, "Eliciting Coordination with Rebates," *Transportation Science*

S. Huang, Q. Wu, "Dynamic Tariff-Subsidy Method for PV and V2G Congestion Management in Distribution Networks," *IEEE Transactions on Smart Grid*

T. A. Taylor, "Supply Chain Coordination Under Channel Rebates with Sales Effort Effects," *Tech. Rep.*

# Carrots vs Sticks

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## Taxes/Tolls

(+)

Added cost to users



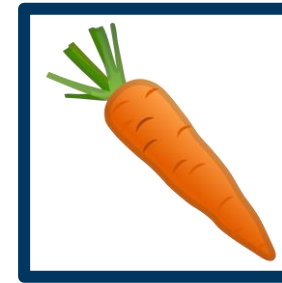
Tolls



## Subsidies/Rebates

(-)

Reduced cost to users



Subsidies

Both are viable methods of influencing users in many settings  
Both can be implemented with similar technology/infrastructure  
Both can be monetarily feasible (fees vs reimbursements to up front cost)

Q?: What are the capabilities of each incentive?

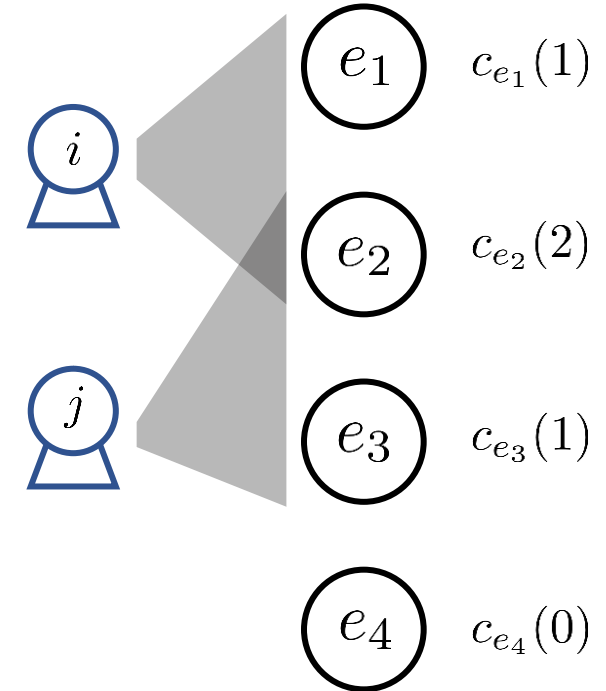
# Model

- Congestion Game  $G$

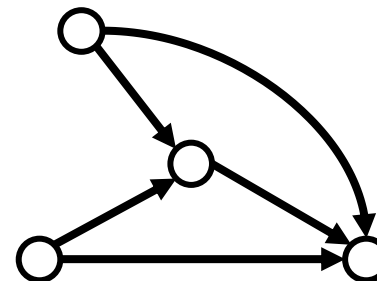
- Resources  $\mathcal{E} = \{1, \dots, E\}$
- Agents  $N = \{1, \dots, n\}$
- Actions  $a_i \in \mathcal{A}_i \subset 2^{\mathcal{E}}$ 
  - Allocation  $a = (a_1, \dots, a_n) \in \mathcal{A}$
- Cost functions  $c_e(|a|_e) \geq 0$
- Total Cost

$$C(a) = \sum_{e \in \mathcal{E}} |a|_e c_e(|a|_e)$$

- Optimal allocation  $a^{\text{opt}} \in \arg \min_{a \in \mathcal{A}} C(a)$



Eg:  
Network  
Congestion



# Selfish Decision Making

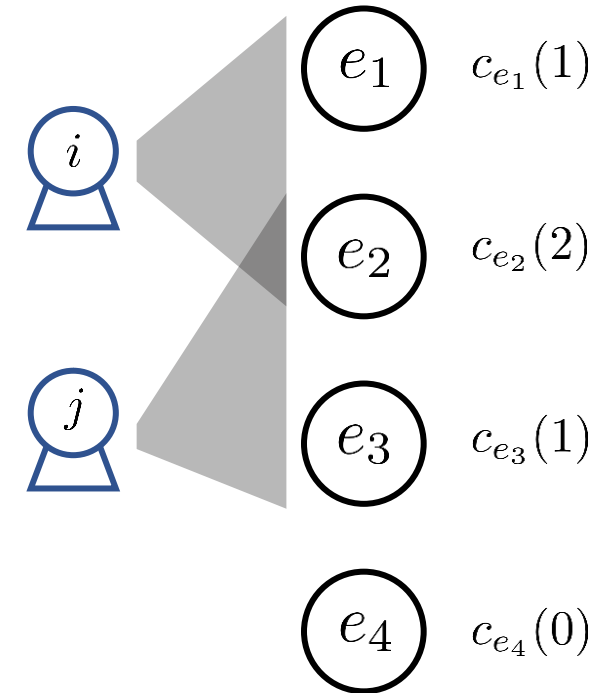
What happens when users choose their own routes?

Cost of user  $i$

$$J_i(a_i, a_{-i}) = \sum_{e \in a_i} c_e(|a|_e)$$

Equilibrium: Nash equilibrium  $a^{\text{Ne}}$

$$a_i^{\text{Ne}} \in \arg \min_{a_i \in \mathcal{A}_i} \sum_{e \in a_i} c_e(|a|_e^{\text{Ne}}) \quad \forall i \in N$$



Price of Anarchy

$$\text{PoA}(G) := \frac{\max_{a^{\text{Ne}} \in \text{NE}(G)} C(a^{\text{Ne}})}{C(a^{\text{opt}})} \geq 1$$

How to reduce this inefficiency?



Incentive Mechanism:  $T(c_e) = \tau_e$

# Selfish Decision Making

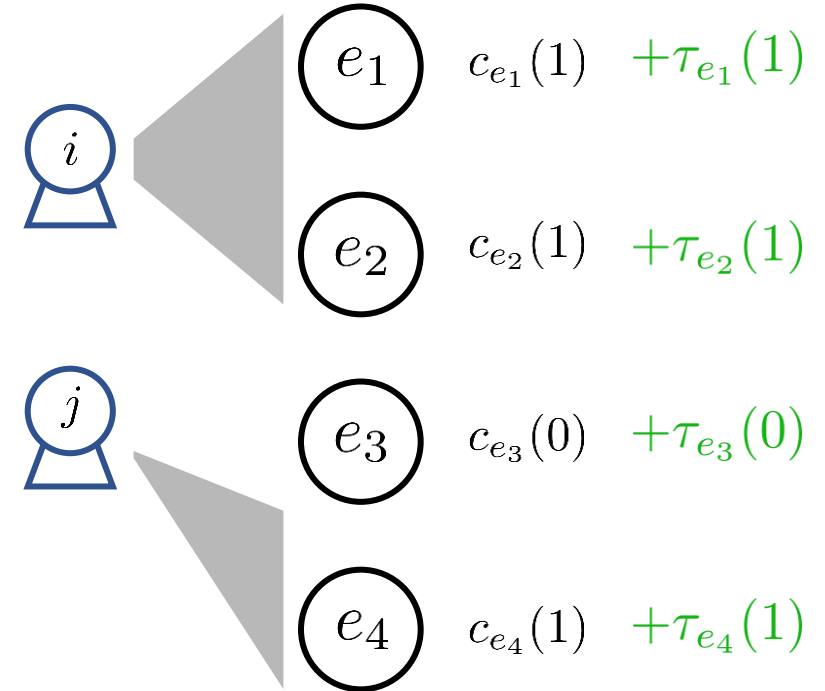
What happens when users choose their own routes?

Cost of user  $i$

$$J_i(a_i, a_{-i}) = \sum_{e \in a_i} c_e(|a|_e) + \tau_e(|a|_e)$$

Equilibrium: Nash equilibrium  $a^{\text{Ne}}$

$$a_i^{\text{Ne}} \in \arg \min_{a_i \in \mathcal{A}_i} \sum_{e \in a_i} (c_e(|a|_e^{\text{Ne}}) + \tau_e(|a|_e^{\text{Ne}})) \quad \forall i \in N$$



Price of Anarchy

$$\text{PoA}(G, T) := \frac{\max_{a^{\text{Ne}} \in \text{NE}(G, T)} C(a^{\text{Ne}})}{C(a^{\text{opt}})} \geq 1$$

How to reduce this inefficiency?



Incentive Mechanism:

$$T(c_e) = \tau_e$$

# Incentives: Taxes & Subsidies

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## Taxes

Tax function:

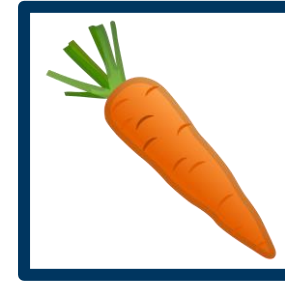
$$\tau_e^+(x) \geq 0 \quad \forall x \geq 0$$

Taxation mechanism:

$$T^+(c_e) = \tau_e^+ \quad \text{Only assigns tolls}$$

Optimal taxation mechanism:

$$T^{\text{opt}+} \in \arg \min_{T^+} \text{PoA}(G, T^+)$$



## Subsidies

Subsidy function:

$$\tau_e^-(x) \leq 0 \quad \forall x \geq 0$$

Subsidy mechanism:

$$T^-(c_e) = \tau_e^- \quad \text{Only assigns subsidies}$$

Optimal subsidy mechanism:

$$T^{\text{opt}-} \in \arg \min_{T^-} \text{PoA}(G, T^-)$$

Tolls  
 $\text{PoA}(G, T^{\text{opt}+})$



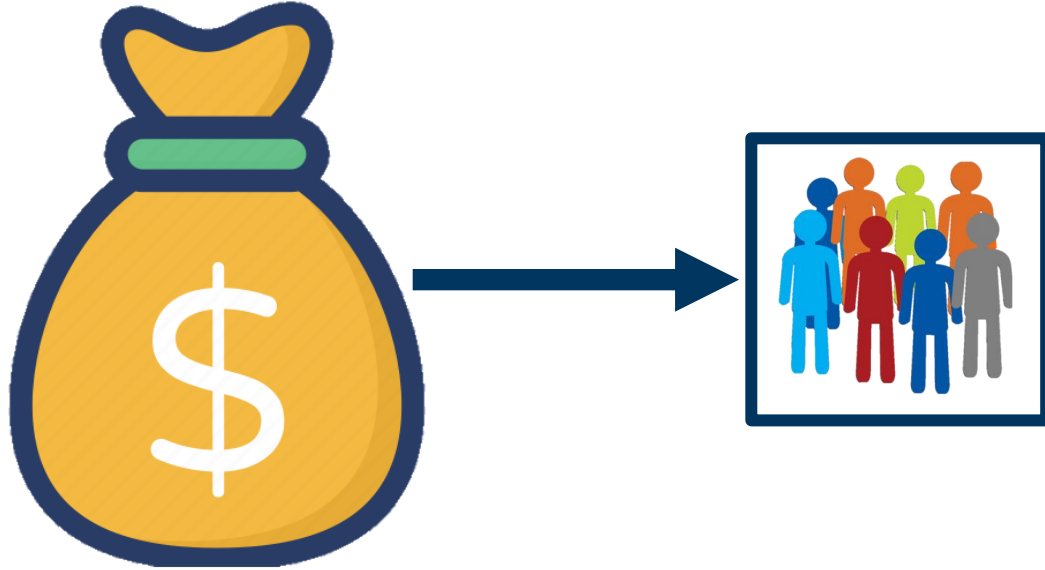
Subsidies  
 $\text{PoA}(G, T^{\text{opt}-})$



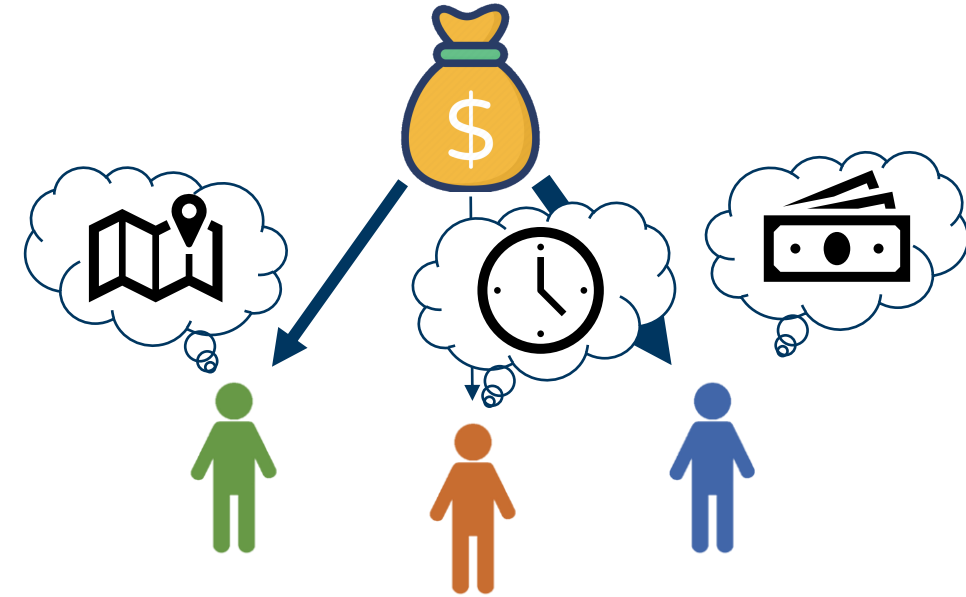
# Challenges

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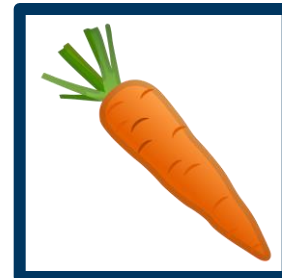
Budgetary constraints



User heterogeneity



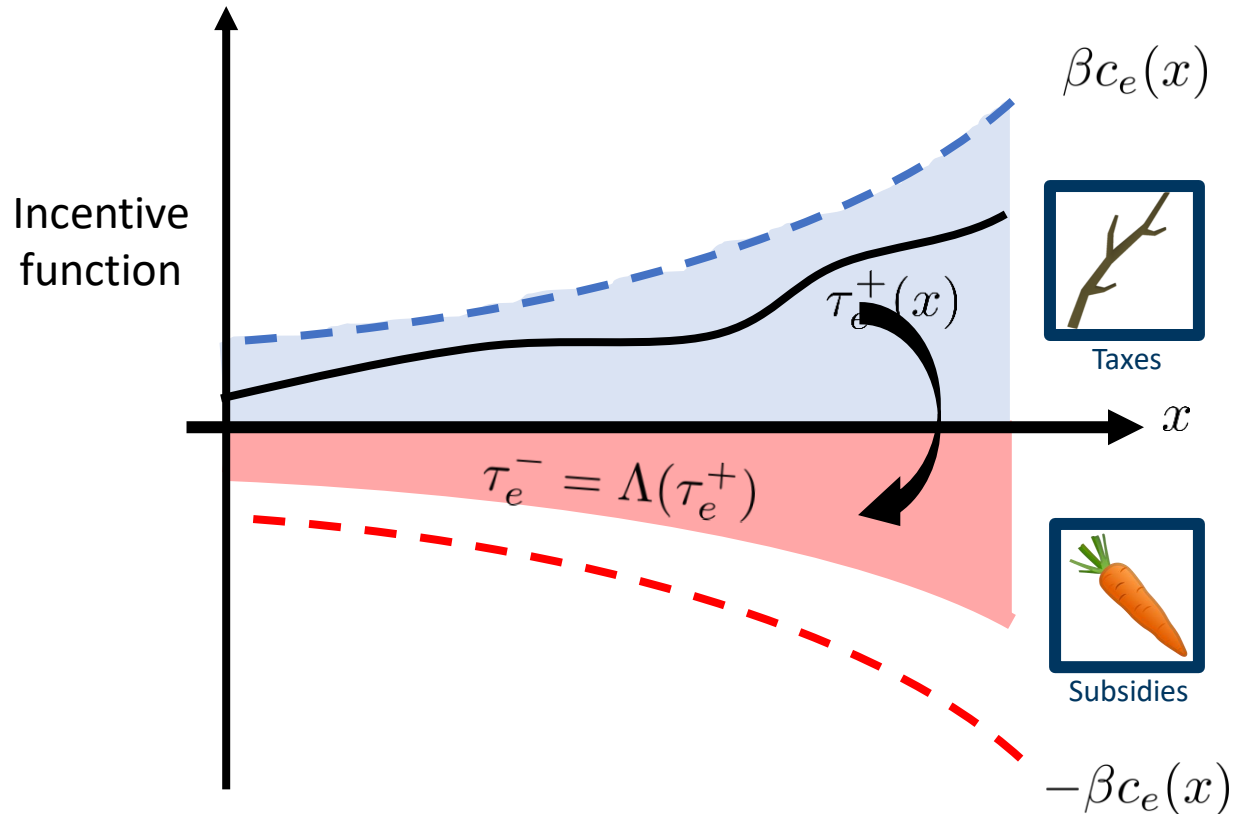
Taxes



Subsidies

# Budgetary Constraints

Added Constraint:  $|\tau_e(x)| \leq \beta c_e(x) \quad \forall x \geq 0$



## Theorem 1

For a congestion games  $G$ , under bounding factor  $\beta \geq 0$ ,

$$\text{PoA}(G, T^{\text{opt}+}(\beta)) \geq \text{PoA}(G, T^{\text{opt}-}(\beta)) \geq 1.$$

Additionally, if the budget constraint is active for every optimal incentive, the inequalities are strict.

**Smaller subsidies can outperform larger taxes.**

# User Heterogeneity

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Each user has unknown price-sensitivity  $s_i \in [S_L, S_U]$

$$J_i(a_i, a_{-i}) = \sum_{e \in a_i} c_e(|a|_e) + s_i \tau_e(|a|_e)$$

$0 < S_L/S_U \leq 1$   
is a measure of  
uncertainty

Robust performance guarantee:

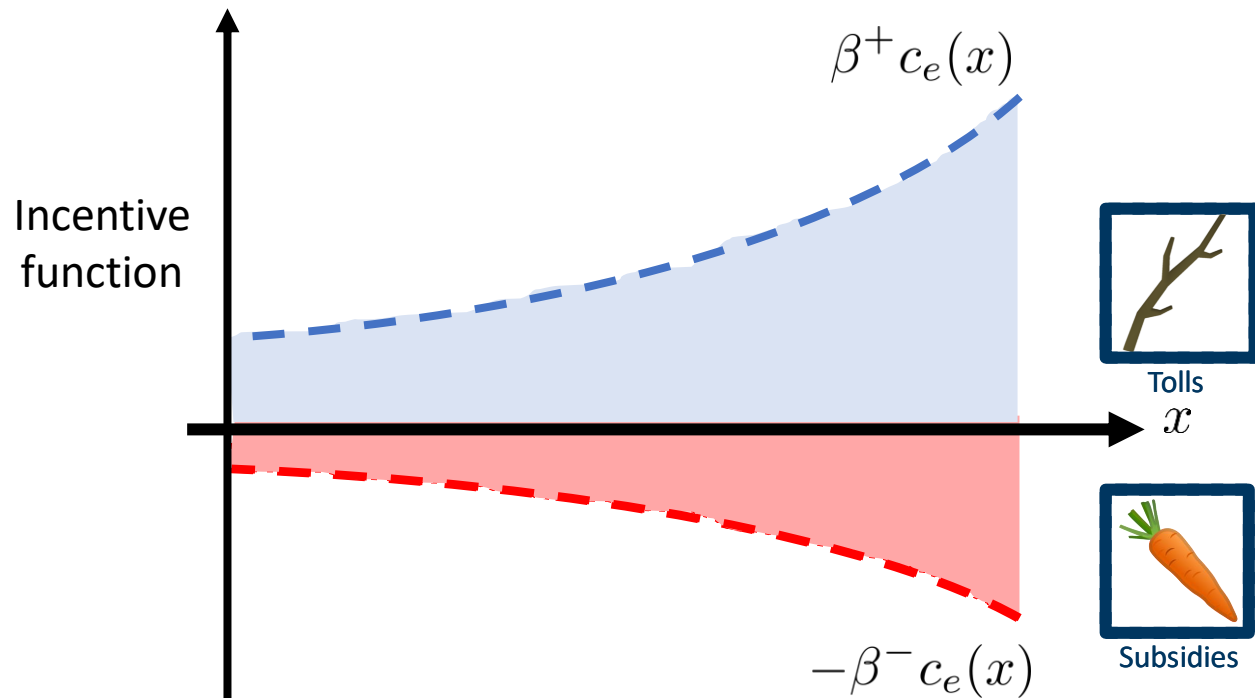
$$\text{PoA}(G, \mathcal{S}, T) = \sup_{s \in \mathcal{S}} \frac{\max_{a^{\text{Ne}} \in \text{NE}(G, s, T)} C(a^{\text{Ne}})}{C(a^{\text{opt}})}$$

Q?: How do incentives perform with  
user heterogeneity?

# Budgetary Constraints & User Heterogeneity

Start with *nominally equivalent* bounded subsidies and tolls, i.e.,

$$\text{PoA}(G, T^{\text{opt}^+}(\beta^+)) = \text{PoA}(G, T^{\text{opt}^-}(\beta^-)) \text{ when users are homogeneous.}$$



As user become heterogeneous:

## Theorem 2

For a congestion games  $G$ , under bounding factors  $\beta^+$ ,  $\beta^-$  respectively, with possible price-sensitivity distributions  $\mathcal{S}$ ,

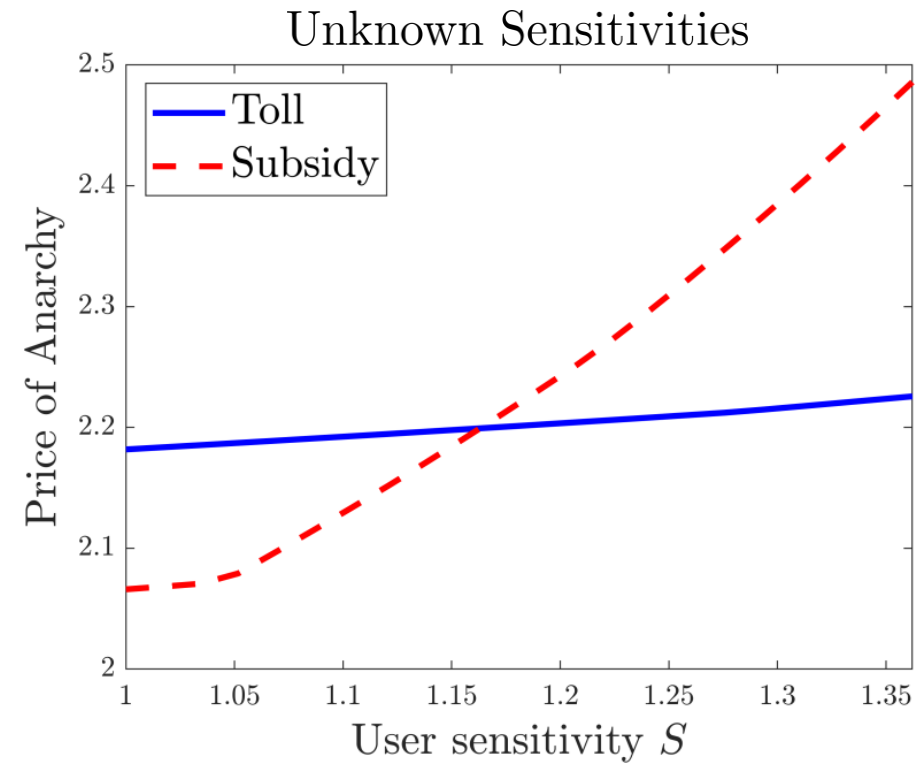
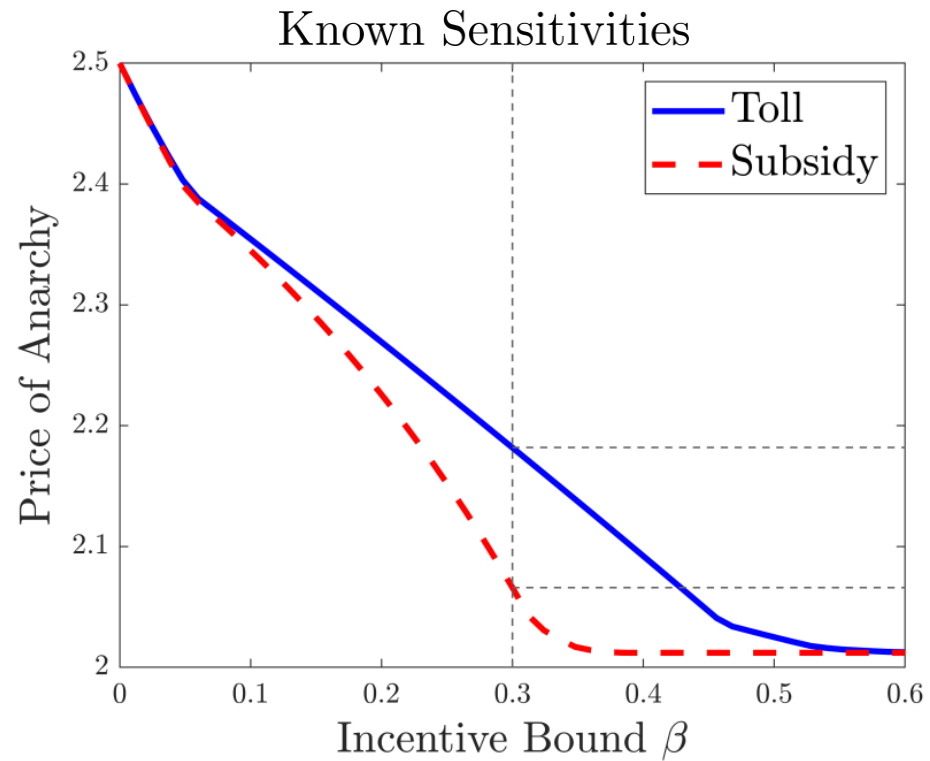
$$\begin{aligned} \text{PoA}(G, \mathcal{S}, T^{\text{opt}^-}(\beta^-, \mathcal{S})) \\ \geq \text{PoA}(G, \mathcal{S}, T^{\text{opt}^+}(\beta^+, \mathcal{S})) \geq 1. \end{aligned}$$

Additionally, if  $G$  is responsive to user heterogeneity, the inequalities are strict.

Performance of *subsidies is less robust* to player heterogeneity than taxes.

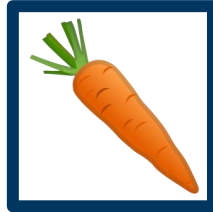
# Computational Example

Price of anarchy bound over affine congestion games.



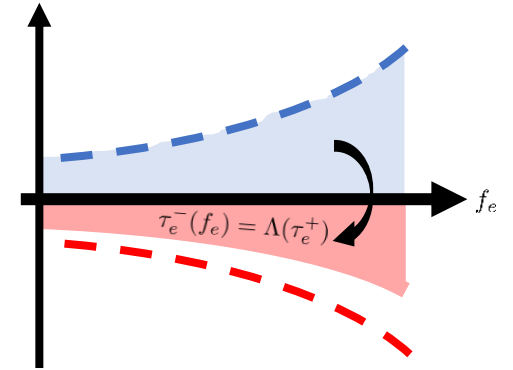
# Conclusion

- Low heterogeneity



Subsidies

Thm. 1

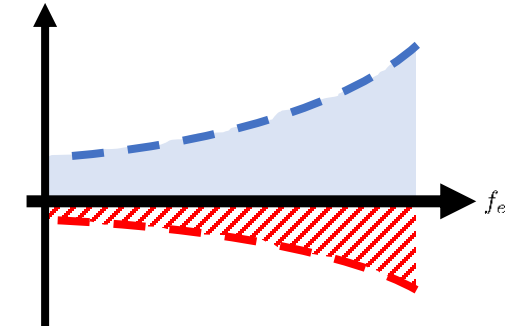


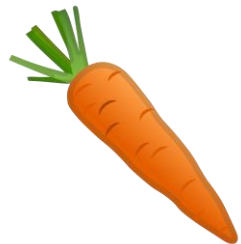
- High heterogeneity



Taxes

Thm. 2





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